

# **Introduction to AI**

**Lecture 15**

## **Inference in First-Order Logic**

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# Inference in First-Order Logic

Inference in First-Order Logic is used to deduce new facts or sentences from existing sentences. Before understanding the FOL inference rule, let's understand some basic terminologies used in FOL.

## Substitution:

Substitution is a fundamental operation performed on terms and formulas. It occurs in all inference systems in first-order logic. The substitution is complex in the presence of quantifiers in FOL. If we write  $F[a/x]$ , so it refers to substitute a constant "a" in place of variable "x".

**Note:** First-order logic is capable of expressing facts about some or all objects in the universe.

## Equality:

First-Order logic does not only use predicate and terms for making atomic sentences but also uses another way, which is equality in FOL. For this, we can use **equality symbols**. **Example: Brother (John) = Smith.**

the object referred by the **Brother (John)** is similar to the object referred by **Smith**. The equality symbol can also be used with negation to represent that two terms are not the same objects.

**Example:**  $\neg(x=y)$  which is equivalent to  $x \neq y$ .

# FOL inference rules for quantifier

- As propositional logic we also have inference rules in first-order logic, so the following are some basic inference rules in FOL:
  - **Universal Generalization**
  - **Universal Instantiation**
  - **Existential Instantiation**
  - **Existential introduction**

## 1. **Universal Generalization:**

• Universal generalization is a valid inference rule that states that if premise  $P(c)$  is true for any arbitrary element  $c$  in the universe of discourse, then we can have a conclusion as  $\forall x P(x)$  are true.

• It can be represented as: 
$$\frac{P(c)}{\forall x P(x)}$$

• This rule can be used if we want to show that every element has a similar property.

• In this rule,  $x$  must not appear as a free variable.

**Example:** Let's represent,  $P(c)$ : "A byte contains 8 bits", so for  $\forall x P(x)$  "All bytes contain 8 bits.", it will also be true.

# FOL inference rules for quantifier

## 2. Universal Instantiation:

- Universal instantiation also called universal elimination is a valid inference rule. It can be applied multiple times to add new sentences.
- The new KB is logically equivalent to the previous KB.
- As per UI, **we can infer any sentence obtained by substituting a ground term for the variable.**
- The UI rule states that we can infer any sentence  $P(c)$  by substituting a ground term  $c$  (a constant within domain  $x$ ) from  $\forall x P(x)$  **for any object in the universe of discourse.**
- It can be represented as: 
$$\frac{\forall x P(x)}{P(c)}$$

### Example:1.

IF "Every person likes ice cream" $\Rightarrow \forall x P(x)$  so we can infer that "John likes ice cream"  $\Rightarrow P(c)$

# FOL inference rules for quantifier

## Example: 2.

Let's take a famous example,

"All kings who are greedy are Evil." So let our knowledge base contains this detail as in the form of FOL:

**$\forall x \text{ king}(x) \wedge \text{greedy}(x) \rightarrow \text{Evil}(x),$**

So from this information, we can infer any of the following statements using Universal Instantiation:

**$\text{King}(\text{John}) \wedge \text{Greedy}(\text{John}) \rightarrow \text{Evil}(\text{John}),$**

**$\text{King}(\text{Richard}) \wedge \text{Greedy}(\text{Richard}) \rightarrow \text{Evil}(\text{Richard}),$**

**$\text{King}(\text{Father}(\text{John})) \wedge \text{Greedy}(\text{Father}(\text{John})) \rightarrow \text{Evil}(\text{Father}(\text{John}))$**

# FOL inference rules for quantifier

## 3. Existential Instantiation:

- Existential instantiation is also called Existential Elimination, which is a valid inference rule in first-order logic.
- It can be applied only once to replace the existential sentence.
- The new KB is not logically equivalent to the old KB, but it will be satisfiable if the old KB was satisfiable.
- What this rule says is that if P holds for some element of the universe, then we can give that element a name such as c (or x, y, a etc). When selecting symbols, one must select them one at a time and must not use a symbol that has already been selected within the same reasoning/proof.
- The restriction with this rule is that c used in the rule must be a new term for which P(c) is true.

It can be represented as: 
$$\frac{\exists x P(x)}{P(c)}$$

### Example:

If you get a 95 on the final exam, then you get an O for the course. Someone, called c, gets 95 on the final exam. Therefore c gets an O. This argument uses Existential.

# FOL inference rules for quantifier

## 4. Existential introduction

- An existential introduction is also known as an existential generalization, which is a valid inference rule in first-order logic.
- This rule states that if there is some element  $c$  in the universe of discourse that has a property  $P$ , then we can infer that there exists something in the universe that has the property  $P$ .

• It can be represented as: 
$$\frac{P(c)}{\exists xP(x)}$$

• **Example: Let's say that,**

"X got good marks in English."

"Therefore, someone got good marks in English."

# Generalized Modus Ponens Rule

For the inference process in FOL, we have a single inference rule which is called Generalized Modus Ponens. It is a lifted version of Modus ponens.

Generalized Modus Ponens can be summarized as " P implies Q and P is asserted to be true, therefore Q must be True."

According to Modus Ponens, for atomic sentences  $p_i, p_i', q$ . Where there is a substitution  $\theta$  such that  $SUBST(\theta, p_i') = SUBST(\theta, p_i)$ , it can be represented as:

$$\frac{p_1', p_2', \dots, p_n', (p_1 \wedge p_2 \wedge \dots \wedge p_n \Rightarrow q)}{SUBST(\theta, q)}$$

## Example:

**We will use this rule for Kings are evil, so we will find some x such that x is king, and x is greedy so we can infer that x is evil.**

Here let say,  $p_1'$  is king(John)       $p_1$  is king(x)  
 $p_2'$  is Greedy(y)                       $p_2$  is Greedy(x)  
 $\theta$  is {x/John, y/John}                   $q$  is evil(x)  
 $SUBST(\theta, q)$  is evil(John).

# Exercise

- From Likes(Jerry, IceCream) it seems reasonable to infer  $\exists x \text{ Likes}(x, \text{IceCream})$ . Write down a general inference rule, **Existential Introduction**, that sanctions this inference. State carefully the conditions that must be satisfied by the variables and terms involved.
- Suppose a knowledge base contains just one sentence,  $\exists x \text{ AsHighAs}(x, \text{Everest})$ . Which of the following are legitimate results of applying Existential Instantiation?
  - a.  $\text{AsHighAs}(\text{Everest}, \text{Everest})$ .
  - b.  $\text{AsHighAs}(\text{Kilimanjaro}, \text{Everest})$ .
  - c.  $\text{AsHighAs}(\text{Kilimanjaro}, \text{Everest}) \wedge \text{AsHighAs}(\text{BenNevis}, \text{Everest})$
  - (after two applications).



# Exercise

- Write down logical representations for the following sentences, suitable for use with
  - Generalized Modus Ponens:
    - a. Horses, cows, and pigs are mammals.
    - b. An offspring of a horse is a horse.
    - c. Bluebeard is a horse.
    - d. Bluebeard is Charlie's parent.
    - e. Offspring and parent are inverse relations.
    - f. Every mammal has a parent.